MobiT: A Distributed and Congestion-Resilient Trajectory Based Routing Algorithm for Vehicular Delay Tolerant Networks

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ABSTRACT
Packet routing is important for Vehicular Delay Tolerant Networks (VDTNs). Opportunistic routing algorithms based on historical records are insufficiently accurate in forwarder selection due to movement randomness of vehicles. Trajectory-based routing algorithms tackle vehicle movement randomness but cannot be directly used in VDTNs due to the dependence on APs. In this paper, we develop a distributed trajectory-based routing algorithm (called MobiT) for VDTNs. This non-trivial task faces three challenges. First, vehicle trajectories must be sufficiently collected. Second, the trajectories cannot be updated frequently due to limited resources of the repository nodes. Third, achieving high routing performance even with partially collected trajectories. Our real trace study lays the foundation of the design of MobiT. Taking advantage of different roles of vehicles, MobiT uses service vehicles that move in wide areas to collect vehicle trajectories, and rely on the service vehicles and roadside units (called schedulers) for routing scheduling. By using regular temporal congestion state of road segments, MobiT schedules the packet to arrive at a roadside unit prior to the destination vehicle to improve routing performance. Further, MobiT leverages vehicles’ long-term mobility patterns to assist routing. Extensive trace-driven and real experiments show the effectiveness and efficiency of MobiT.

CCS CONCEPTS
- Networks → Network control algorithms; Cyber-physical networks;

KEYWORDS
VDTN, distributed routing, vehicle trajectory

1 INTRODUCTION
In recent few years, many research efforts have been devoted to Vehicular Delay Tolerant Networks (VDTNs) [1, 8]. VDTNs can alleviate bandwidth burden on networks and serve areas with sparse infrastructures. In such vehicular networks with sparse connection, packet delivery between vehicles is important for many purposes. For example, a vehicle needs to report a traffic accident to a police vehicle far from the crash site, or a vehicle using a mobile social network wants to share a newsfeed with its friend vehicle miles away. However, due to the high mobility of vehicles and disconnected nature of VDTNs, efficient packet delivery is non-trivial.

Previous opportunistic routing algorithms [7, 8, 21, 22] define different utilities (e.g., meeting probabilities) and forward the packet to vehicles or Roadside Units (RSUs) that have larger utilities with the destination vehicle. However, these algorithms use vehicles’ historical meeting records to schedule packet forwarding, which has been proven insufficiently accurate [27] due to movement randomness of some vehicles.

Determining packet forwarder based on vehicles’ trajectories is effective in handling movement randomness [10–12, 24]. In the trajectory-based routing algorithms, vehicles repeatedly report trajectories to Access Points (APs) sparsely located along roads. A central server then uses these shared trajectories to schedule forwarders to carry the packet to the destination vehicle in its driving route. However, these algorithms cannot be directly used in VDTNs due to the dependence on APs.

In this paper, we design MobiT, which derives vehicle Mobility from Trajectories for routing. First, MobiT uses service vehicles to collect vehicle trajectories, and relies on the service vehicles and roadside units (RSU) (both are called schedulers) for determining routing path. Second, MobiT only requires each participating vehicle to report its trajectory to a scheduler when it starts moving (called initial trajectory) rather than repeated reporting. To tackle outdated trajectory, MobiT considers the temporal change of road congestion state when using the initial trajectories for vehicle movement prediction. Third, MobiT exploits both short-term mobility (i.e., trajectory) and long-term mobility (i.e., road/area visiting pattern) as complementary approaches. When determining routing path, MobiT schedules the packet to arrive at an RSU prior to the destination vehicle, which generates higher performance than scheduling direct meeting as in the previous trajectory-based algorithms. When a routing path cannot be found, MobiT finds a path to let the packet approach the destination vehicle. If the trajectory of the destination vehicle is unavailable, MobiT uses the vehicle’s long-term mobility to forward the packet.

MobiT can also overcome some problems in the previous centralized trajectory-based methods. First, due to fluctuating road
traffic, it is very difficult to schedule an exact meeting with the destination vehicle regardless of the powerful capacity of the central server. Second, vehicles sometimes rely vehicular communication to maintain contact with the APs. The trajectories in the central server may be outdated due to vehicles’ intermittent connection to the APs and possible packet loss in communication. Third, the selection of forwarders does not consider the change of road traffic at different times and road segments. To our best knowledge, this work is the first that realizes efficient distributed trajectory-based routing algorithm in VDTNs. The remainder of the paper is organized as follows. Section 2 presents literature overview. Section 3 presents the design of MobiT. Section 4 presents the performance evaluation. Section 5 concludes this paper with future directions.

2 RELATED WORK

Opportunistic routing algorithms. These algorithms extract utilities from vehicles’ historical records. The packet is forwarded at the direction maximizing the utility. SADV [7] lets packets wait at intersections until the path with minimum delay is available. Ishihara et al. [8] schedules packet delivery by the packet’s aggregated demand and the historical condition of neighboring vehicles having the same packet. EBT [21] utilizes users’ previous encounters to construct a relation graph for packet forwarding. Tie et al. [22] proposed the Robust Replication Routing (R3), which unifies mesh, MANET, DTN routing paradigms by predicting the distribution of link delays. Kong et al. [14] proposed a frequency-divided instantaneous neighbors estimation system for vehicular networks. In [19], Schwartz et al. focuses on guidelines for the design of data dissemination in vehicular networks. These works rely on historical information to predict future encounter, but cannot guarantee high efficiency.

Trajectory-based routing algorithms. Several recent works utilizing vehicles’ trajectories in routing have been proposed. Wu et al. [23] found and used the spatio-temporal correlation of vehicle mobility in data delivery. TBD [11], TSF [12] and STDFS [24] use APs to collect vehicle trajectories. Then, the rendezvous position between the destination vehicle and the packet is determined based on accumulated trajectories, and the packet is forwarded to the rendezvous position. Due to the sparsity of APs, ad-hoc network is needed to bridge APs and vehicles. However, these algorithms have drawbacks of susceptibleness to road congestion and possibly outdated trajectories due to inaccessibility to APs, which affect the accuracy of the forwarder selection.

3 SYSTEM DESIGN

We consider a VDTN with \( n \) vehicles denoted by \( N_i (i = 1, 2, \ldots, n) \) and make following assumptions.

(1) Each vehicle is equipped with a Dedicated Short Range Communication (DSRC) device [5]. When two vehicles are within each other’s communication range, an encounter happens.

(2) Each vehicle is equipped with a navigation system, which generates trajectory consisting of future positions and estimated arrival times, and road maps [15, 25].

(3) The area in the VDTN is partitioned into multiple sub-districts with equal number of landmarks. Following [26], we assign a center landmark for each sub-district.

(4) Each intersection is installed with an RSU which uses DSRC for communication [6, 18]. Service vehicle can exchange information with RSUs, while the others can only drop a packet to an RSU.

There are existing works focusing on motivating users to share mobility information [4] and ensuring users’ privacy [16]. We leave the work for MobiT as our future work.

3.1 Representation of Vehicle Mobility

3.1.1 Short-term Mobility and Congestion-considered Update.

At the beginning of a trip, each vehicle generates its initial trajectory. An initial trajectory of vehicle \( N_i \) is \( \langle N_i; [p_1^i, p_2^i, \ldots, p_Q^i]; T_i \rangle \), where \( [p_1^i, p_2^i, \ldots, p_Q^i] \) represents the sequence of \( Q \) positions on the trajectory. \( T_i \) is the starting time of this trajectory. Each vehicle maintains its own short-term mobility information. While service vehicle collects short-term mobility information from every vehicle it meets. If a service vehicle meets another service vehicle or an RSU, they exchange their known mobility information.

Congestion state table: MobiT use road segment, which is the interval between two neighbor intersections, as the basic unit of roads. It has been shown that the travel time on a road can be estimated based on the congestion state of composing road segments [13]. It is also noticeable that urban traffic pattern repeats in daily fashion [9]. Thus, we use the congestion states of roads under different times to assist determining vehicle arrival times. We firstly design the table of delays, which records the travel times of a road’s composing segments under congested and non-congested cases based on historical statistics. Then, for each road, we design a table of binary vectors to describe its road congestion.

For each road segment, it has distinct travel times corresponding to congested and non-congested situations [13]. For example, College Ave has 6 segments as shown in Table 1. If segment 1 is congested, it takes the vehicle 2min to drive through. Otherwise, 50s is needed for driving through. Suppose the segments 1, 2 and 4 are congested, and the other segments are non-congested, the travel time needed to drive through College Ave is 50s + 5min + 6min + 20s + 2min + 10s = 14min20s.

<table>
<thead>
<tr>
<th>Segment ID</th>
<th>0s</th>
<th>5s</th>
<th>1min</th>
<th>2min</th>
<th>5min</th>
<th>10min</th>
</tr>
</thead>
<tbody>
<tr>
<td>Congested (1)</td>
<td>2min</td>
<td>3min</td>
<td>6min</td>
<td>1min</td>
<td>2min</td>
<td>5min</td>
</tr>
<tr>
<td>Otherwise (0)</td>
<td>50s</td>
<td>2min</td>
<td>1min</td>
<td>20s</td>
<td>1min</td>
<td>10s</td>
</tr>
</tbody>
</table>

Then, for each sequence of road segments, we can use a binary vector to depict its congestion states. For example, if the segments 1, 2 and 4 of College Ave are congested, current congestion state of the road is \([0, 1, 1, 0, 1, 0]\). To collect all possible congestion states of a road during different times, the congestion state of the road is sampled by a time unit, say per hour. Thus, for each hour, we have several sampling results representing all the possible congestion states of the road at this time. Finally, we classify these congestion states along with their respective probabilities in ascending order of time, and get the table of congestion states of the road as shown in Table 2. In Table 2, College Ave has several congestion states under
each time interval. Each congestion state has a probability, which measures its appearance frequency among all possible congestion states. For example, during the interval from 00:00 to 01:00, College Ave has the probability of 0.6 to be in congestion state [0, 1, 0, 0, 1, 0] and the probability of 0.4 to be in congestion state [1, 0, 0, 0, 1, 0]. Therefore, the estimated travel time of College Ave during this interval is $0.6 \times 14 \text{min} 20\text{s} + 0.4 \times 7\text{min} 30\text{s} = 11\text{min} 36\text{s}$.

Table 2: Table of College Ave's congestion states.

<table>
<thead>
<tr>
<th>Time interval</th>
<th>Congestion states</th>
</tr>
</thead>
<tbody>
<tr>
<td>00:00−01:00</td>
<td>[0, 1, 0, 0, 1, 0]</td>
</tr>
<tr>
<td></td>
<td>[1, 0, 0, 0, 1, 0]</td>
</tr>
<tr>
<td></td>
<td>[0, 0, 1, 1, 0, 0]</td>
</tr>
<tr>
<td></td>
<td>[0, 0, 0, 1, 1, 0]</td>
</tr>
<tr>
<td></td>
<td>[1, 0, 0, 0, 0, 1]</td>
</tr>
<tr>
<td></td>
<td>[0, 0, 0, 0, 1, 1]</td>
</tr>
</tbody>
</table>

Since road traffic follows certain long-term pattern even under accident and weather change, using historical data to describe the congestion states is reasonable [13, 20]. The table of congestion states and delays for all roads are computed offline. All schedulers are preloaded with the two tables.

**Estimation of travel time and deviation:** To estimate the travel time of a trajectory, a scheduler decomposes it into roads. For each road, the scheduler refers to Table 2 for the road’s current congestion state. Then, the scheduler refers to Table 1 for corresponding delays of the covered segments. By summing the delays from the start of each trajectory, the scheduler estimates the vehicle’s future travel time. For example, suppose a vehicle will drive through College Ave between 00:00 and 01:00. According to Tables 1 and 2, the travel time of College Ave has probability of 0.6 to be 14min20s and probability of 0.4 to be 7min30s. Therefore, the road’s estimated travel time is $\mu = 0.6 \times 14\text{min} 20\text{s} + 0.4 \times 7\text{min} 30\text{s} = 11\text{min} 36\text{s}$, with standard deviation $\sigma = \sqrt{0.6 \times (14\text{min} 20\text{s} - 11\text{min} 36\text{s})^2 + 0.4 \times (7\text{min} 30\text{s} - 11\text{min} 36\text{s})^2} = 3\text{min} 20\text{s}$.

3.1.2 Long-term Mobility. In this section, we introduce long-term mobility, namely routine and vehicular friendship.

**Routine:** Vehicles’ long-term mobility has regularity [27], which is reflected as certain roads that are frequently driven by the vehicle at specific times. For example, people usually take the same routine routes to commute between home and work place. Moreover, vehicles tend to repeat their routine routes on a daily basis. For example, people regularly drive from home to work place at around 8:10 every morning. Therefore, we depict the routines of a vehicle, say $N_i$, as shown in Table 3.

Table 3: Table of a vehicle’s routines.

<table>
<thead>
<tr>
<th>Pool</th>
<th>Route</th>
<th>$T_s$</th>
<th>$T_e$</th>
</tr>
</thead>
<tbody>
<tr>
<td>0.6</td>
<td>$\langle p^0_{i1}, \ldots, p^m_{i1} \rangle$</td>
<td>09:10 $\sim$ 08:20</td>
<td>09:30 $\sim$ 08:45</td>
</tr>
<tr>
<td>0.2</td>
<td>$\langle p^0_{i2}, \ldots, p^m_{i2} \rangle$</td>
<td>10:00 $\sim$ 12:20</td>
<td>13:30 $\sim$ 13:45</td>
</tr>
</tbody>
</table>

In Table 3, each row represents a routine of $N_i$. Route represents the series of positions a routine covers. $T_s$ is the start time range of the routine. $T_e$ is the end time range of the routine. $T_s$ and $T_e$ are determined from $N_i$’s historical records. The probability indicates the likelihood that the vehicle will follow the routine and is calculated as $Pr_{R_i}(N_i) = m_i / M$, where $m_i$ is the number of occurrences that $N_i$ followed routine $R_i$ as a trajectory; $M$ is the total number of trajectories that $N_i$ once drove during a time period, say 30 days. The routine table stores the routines with the sum of probabilities larger than and closest to 80%. This threshold can be adjusted based on system constraints.

**Friendship:** People living in the same area are likely to follow similar routines. For example, suppose Alice and Bob live in the same suburban community, every morning they drive the same highway to downtown. Given a packet targeting at Alice, the mobility information of Bob may be helpful. Based on such observation, MobiT measures the relationship between vehicles in terms of their similarity on routine.

**Overlapping of routines:** We define two vehicles are friends if the ratios of their similar routines in their respective overall routines are higher than a threshold, say $\alpha$. For example, suppose vehicle $N_i$ has routines: $R_1$, $R_2$ and $R_3$, vehicle $N_j$ has routines: $R_2$ and $R_3$, and $\alpha = 0.5$. Since $R_2$, $R_3$ are the similar routines of $N_i$ and $N_j$, and they take up 66.7% and 100% of the total routines of $N_i$ and $N_j$, respectively, which are higher than 0.5, these two vehicles are friends.

Since the similar routines of two vehicles will not be completely identical, we use spatiotemporal overlap to measure their similarity. Given two routines, say $R_1$ of $N_i$ and $R_2$ of $N_j$, $R_1$ covers positions: $r_1 = \{ p_{i1}^0(0), \ldots, p_{i1}^m(m) \}$, start time range $T_{s1}$ and end time range $T_{e1}$, while $R_2$ covers positions: $r_2 = \{ p_{j2}^0(0), \ldots, p_{j2}^m(m) \}$, start time range $T_{s2}$ and end time range $T_{e2}$. We use $\tilde{T}_s$ and $\tilde{T}_e$ to denote the mean of $T_s$ and $T_e$, respectively. Then, $R_1$ and $R_2$ are similar if:

$$\left|\frac{T_{s1} - T_{e1}}{\tilde{T}_s}\right| < \tau_t$$

$$\left|\frac{T_{s2} - T_{e2}}{\tilde{T}_e}\right| < \tau_t$$

$$\frac{|p_1 \cap p_2|}{|p_1 \cup p_2|} > \gamma_s$$

where $\tau_t$ is the threshold bounding the temporal deviation of start times and end times. $\gamma_s$ is the threshold bounding the spatial deviation of the positions. The thresholds are determined based on the traffic flow of specific scenes. For metropolitan cities, relatively high deviation should be tolerated. Therefore, we set $\alpha = 0.5$, $\tau_t = 15 \text{min}$ and $\gamma_s = 0.6$.

In MobiT, routine extraction and friendship determination are conducted by service vehicles and RSUs since they have the bulk of vehicles’ short-term mobility information. A representative friend list is as shown in Table 4.

Table 4: Table of friends.

<table>
<thead>
<tr>
<th>Vehicle ID</th>
<th>Friends</th>
</tr>
</thead>
<tbody>
<tr>
<td>N1</td>
<td>N1(10:21) N1(10:24) N1(09:12)</td>
</tr>
<tr>
<td>N2</td>
<td>N2(09:41) N2(10:24)</td>
</tr>
</tbody>
</table>

3.2 Routing Process

MobiT aims to deliver the packet to the encounter position prior to the destination vehicle, the packet then waits at a nearby RSU for the destination vehicle. In MobiT, service vehicles and RSUs schedule the forwarding of packets since they have collected vehicles’ mobility information. The scheduling of routing can be summarized to three cases. 1) When the destination vehicle’s short-term mobility information is available, the packet will be forwarded to the destination vehicle’s future position (or nearby position) along a trajectory-based routing path that leads to the shortest delay; 2) When the scheduler only has the destination vehicle’s long-term mobility information, the packet will be forwarded to the destination vehicle’s or its friend’s routine; 3) When no mobility information of the destination vehicle is available, the packet will be forwarded to service vehicles, aiming to increase its probability of finding more useful information.
3.2.3 Routing without Mobility Information. It is likely the scheduler doesn’t have any useful mobility information. If the scheduler is a service vehicle, it will keep the packet. If the scheduler is a RSU, it will first keep the packet and then transfer it to the service vehicle passing by.

Because each scheduler stores partial mobility information of vehicles in the system, a routing path generated by a scheduler may not be the best routing path. Therefore, whenever a node carrying a packet encounters a scheduler, it requests the scheduler to update the routing path if a chain with shorter delivery delay is found. In the routing process, if the packet misses the next forwarder, it requests nearby scheduler to launch a new round of routing.

4 PERFORMANCE EVALUATION

We used the Rome [2] and the San Francisco [17] traces, which last for 30 days, for evaluation. The Rome trace has has 315 taxis and 4638 landmarks, and the San Francisco trace has 536 taxis and 2508 landmarks. We develop a trace-based simulation environment which is driven by each vehicle’s movement event [3, 4].

The collection of congestion table and delay table was finished offline. For Rome and San Francisco, the threshold speeds to determine congestion are 20MPH and 30MPH, respectively [2, 17]. The congestion state of each road segment was sampled per hour. For both traces, we set the initial period to 7 days, during which service vehicles collected and disseminated mobility information. Meanwhile, service vehicles and RSUs extracted vehicles’ routines as in Section 3.1.2, and determined friendship between vehicles by Equation (1) and (2) with $\alpha_f = 0.5$, $\tau_f = 15\text{min}$ and $\gamma_s = 0.6$. Request rate is the number of packets generated every 24 hours in both traces and was set to 40 by default. Packet TTL, which is the valid time of a packet, was set to 24 hours. The TTL for short-term mobility information depends on trip duration.

We compared MobiT with two representative algorithms: the Shared-Trajectory-based Data Forwarding method (STDFS in short) [24], and the Robust Replication Routing (denoted by R3) [22]. STDFS depends on vehicles’ trajectories reported through APs to schedule future meeting position between forwarder and destination vehicle. In R3, vehicles record their historical contact with others. The packet carrier utilizes the historical delays of the vehicles to the destination vehicle to guide packet routing. In simulations, we equipped 2782 and 1504 landmarks with RSUs/APs in Rome and San Francisco, respectively, which is as specified in STDFS [24]. We measured following metrics:

- **Success rate**: The percentage of packets that successfully reach their destination vehicles.
- **Average delay**: The average time (in seconds) used by packets to reach their destination vehicles.
- **Average number of information queries**: The average number of information queries transmitted among nodes.
- **Average vehicle memory usage**: The average number of memory units used by each vehicle.

4.1 Experimental Results

In the experiment, we varied the packet TTL from 18 hours to 33 hours with 3 hours as the step size.

4.1.1 Success Rate. Figure 1(a) and Figure 2(a) show the success rates of the algorithms under different packet TTLs in both traces. In these figures, the success rates follow: MobiT > STDFS > R3. We can see that MobiT always has the highest packet delivery success rate than the other two algorithms under various situations.

The success rates of all algorithms remain nearly constant under different request rates, but increase with the ascending of packet TTL. R3 always has the lowest success rate. This is because vehicles have independent and random movement, vehicles’ historical meeting records with the destination vehicle does not guarantee another meeting with it in future. Thus the selection of forwarder may be mistaken.

In contrast, STDFS has much higher success rate. It is because the packet is always forwarded to a future position of the destination vehicle with certain accuracy. If vehicles’ movement is not influenced by road congestion, or its trajectory in central server is continuously updated without disconnection, the forwarder is highly likely to meet the destination vehicle.
MobiT always achieves the highest success rate. This is because MobiT considers road congestion state in estimating vehicles’ arrival times on trajectory, which makes the estimation more tolerant to traffic change. Also, MobiT is not afraid that trajectory may be outdated because it uses road delay table corresponding to road congestion to dynamically estimate the arrival times of vehicles when scheduling routing. Moreover, MobiT aims to let packet arrive at the meeting position prior to the destination vehicle through considering various kinds of mobility information, which also increases the success rate.

4.1.2 Average Delay. Figure 1(b) and Figure 2(b) show the metric under different packet TTLs in both traces. In Rome, the average delays follow: MobiT < STDFS < R3. While in San Francisco, the average delays follow: MobiT < R3 < STDFS. We can see that MobiT achieves the best performance.

R3 does not know the position of the destination vehicle, so it is likely to select the vehicle that will not meet the destination vehicle as forwarder. Therefore, it has the highest delay. Although STDFS knows the future position of the destination vehicle from its trajectory, it can only forward packet when a complete chain of trajectories connecting the source vehicle and the destination vehicle is available. Also the destination vehicle’s disconnection to APs will make its trajectory outdated, thereby hindering the efficient delivery of the packet. So STDFS ranks the second. MobiT utilizes various kinds of vehicles’ mobility information to help the packet keep approaching the actual activity area of its destination vehicle, so it has the shortest delay.

4.1.3 Average Number of Information Queries. Figure 1(c) and Figure 2(c) show the metric of the algorithms under different packet TTLs in both traces. The average number of information query follow: R3 < MobiT < STDFS. We can see that MobiT achieves less information query overhead than STDFS but more information query overhead than R3.

This is because STDFS requires vehicles to repeatedly report their trajectories to the APs, so it has the highest number of information query. In contrast, MobiT only needs vehicles to report their initial trajectory to schedulers. Therefore, it ranks the second. In R3, information query only happens in the encounter of nodes with suitable delay predictions. Therefore, it ranks the lowest.

4.1.4 Average Memory Usage. Figure 1(d) and Figure 2(d) show the metric of the algorithms under different packet TTLs in both traces. For MobiT, we additionally measured the metric for service vehicle and RSU, which is represented with “Service”. The metric follows: R3 > Service > MobiT > STDFS in Rome, and Service > R3 > MobiT > STDFS in San Francisco. We can see that the memory usage of general vehicle is comparable to the one of STDFS.

Since R3 has duplicated packets, and each vehicle needs to maintain the distribution of path’s historical delays, vehicles in R3 have the highest memory usage. In MobiT, general vehicles only need to maintain short-term mobility information and occasional packets. Therefore, MobiT uses much lower memory than R3. On the other hand, service vehicles and RSUs need to maintain much mobility information and awaiting packets. Therefore, their memory usage is comparable to that of R3. In STDFS, vehicles only need to record their trajectory. In contrast to MobiT, in which a vehicle may need to help forwarding packet even if they are not determined to approach the destination vehicle. STDFS only requires vehicles that can form a complete chain of trajectories between the source vehicle and the destination vehicle. Therefore, STDFS uses the least memory.

5 CONCLUSION

Message delivery is an important function in VDTNs for Intelligent Transportation Systems. Previous opportunistic routing algorithms for VDTNs cannot achieve high success rate and low delay due to insufficiently accurate estimation of vehicles’ future encounter. Previous trajectory-based routing algorithms can overcome this drawback but require APs hence cannot be directly used for decentralized VDTNs. We propose MobiT, a distributed trajectory-based routing algorithm for VDTNs. MobiT aims to let the packet arrive at a RSU prior to the destination vehicle. By taking advantage of travelling features of different vehicles, MobiT uses public service vehicles and RSUs to collect vehicle mobility information in a distributed manner and schedule trajectory-based routing paths to
destination updates, trajectories only need to be reported once and then are updated based on stored road segment congestion state at different times. Extensive trace-driven and real-world experiments show MobiT’s higher efficiency and effectiveness compared with previous routing algorithms. In the future, we plan to further exploit vehicles’ relationship in routing.

ACKNOWLEDGEMENTS

This research was supported in part by U.S. NSF grants ACI-1719397 and CNS-1733596, and Microsoft Research Faculty Fellowship 8300751.

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